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$\mathbf{1}$ Introduction

A Message Translation and Validation -MTV module is a module common in command and control systems It performs the functions of translation and validation between three main representations of messages External Representation -EXR Internal Representation -INR and user algebrasis are transpiration and the messages are transmitted to and from systems outside of the command and control system. Such messages tend to be terse and/or encoded for efficient communication between systems. Internal representation is how messages are stored and manipulated by the command and control system of which the MTV component is a part. The user representation is a format intended for viewing and input by the users of the system Translation must be performed between EXR and INR representations and between INR and USR representations. Validation must be performed on EXR messages and USR messages; INR messages are presumed to be valid

An MTV generator is a program that takes as input a specication of an MTV module and produces the code for that module. The following document gives a requirements specification for MTV modules and thus specifies what an MTV generator must support.

$\overline{2}$ The MTV representations

- - -The External Representation

An external message consists of "bits on a wire"—simply a string of bits. Thus is has no explicit structure; however, it does have implicit structure described by its message format. According to the format, an external message is subdivided into fields of bits or bytes. Fields may be fixed length or variable length. When variable length, they must be terminated with a delimiter. The contents of a field is an encoding for either symbolic data, numeric data, or text.

For example the following external message consists of a reporting location -loc a Julian date and time, and a text message. The reporting location's value is a code for a particular Air Force Base The date and time is a binary encoding of January at am The text is a variable length byte-field, terminated with a period, and is ASCII for "Hi".

The specification for an EXR is usually part of the requirements levied upon the system that is employing MTV. Typically, these specifications are informal, textual descriptions of the EXR An example informal EXR description is presented in gure This is a rather contrived example, pieced together from several actual message formats, but it serves to demonstrate most of the message features that an MTV system is required to handle Omitted for brevity in this example are the details of the bit layout for the first two fields.

-The Internal Representation

The internal representation is a datatype in the language that the MTV instance is written in In this case, that will be Ada. However, what the MTV generator will accept for an INR will be a subset of Ada datatypes. In particular, the INR can consist of integers, strings, enumerations, records, variant records, arrays and linked lists. The INR for a given message format is defined by the designer of the system employing MTV

An INR is assumed to represent a valid message. For INR messages created by MTV, this is assured by the translation and validation process The application employing MTV is held responsible for the validity of any INR message that it creates

An INR corresponding to the previous example is presented in figure 2. Notice that the elements of the EXR description correspond roughly to the Ada datatype. There are, however, significant differences. For example, the Detection Confidence and Probability of Detect fields, which are separate in the EXR description are intermingled in a single array in the INR.

-The User Representation

The user representation is a user-readable character string. The intent of the USR is for presenting messages to the user in a readable format and as a representation for user input of messages A USR should be easily understood by a user familiar with the message format

Note: if no Contact Data is available, each element in the Detection Confidence field must indicate None

Figure External message format example

```
type Reporting_Location_Type is (Andrews_AFB,
                                              PetersonAFB
                                              Wright_Patterson_AFB);
subtype Julian_Day_Type is Integer range 1 .. 366;
subtype is the Hourst range of \mathcal{L}_1 . In the following range \mathcal{L}_2 , and the following range \mathcal{L}_1subtype Minutes and \mathbf{r} and \mathbf{r} is \mathbf{r} range range
type Julian_Date_Time_Record_Type is record
   JulianDay-
 JulianDayType
   \blacksquare Hour-Hour-Type \blacksquareMinute-
 MinuteType
end record;
the message state in the message of the state of the state of \mathcal{T}_\mathcal{T}type Track_Info_Discr_Type is (Altitude_Discr, Confidence_Discr);
subtype Trackaltitude Type is Integer range field at the p
type Confidence_Type is (High, Middle, Low, None);
\blacksquare type Tracking is recorded to \blacksquare . The correction is recorded to \blacksquarecase Discr is
  when Altitude_Discr =>
      Altitude-
 TrackAltitudeType
  when Confidence_Discr =>
      Confidence-
 ConfidenceType
  end case
end record;
```
Figure 2: Internal message format example

```
type Contact_Data_Type;
type Contact_Data_Link_Type is access Contact_Data_Type;
subtype Track_Number_Type is Integer range 1 .. 9999;
type Force_Code_Type is (Hostile, Unknown, Friendly);
subtype contact and contact in the subtype sharper range of the state \mathcal{L}_\mathbf{A}subtype \blacksquare is \blacksquare is integrated the following range \blacksquaretype Contact_Data_Type is record
    the contract of the contract of
    recommendation in the contract of the code-type of the code-type of the code-type of the code-type of the code-
    — Contact Latitude Contact Latitude 1999, International Contact Latitude 2012
    Longitude-
 ContactLongitudeType
    next-contact the contact of the contact of
end record;
subtype Barrier in the subtype is \mathbf{H} = \mathbf{H} \mathbf{H} . The subtype is \mathbf{H} = \mathbf{H} \mathbf{H}type Barrier_Segment_Type is record
    DetectionConfidence-
 ConfidenceType
    red and the contraction of the c
end record;
type Barrier_Array_Type is array (1 .. 8) of Barrier_Segment_Type;
type Example_Message_Type is record
    R reporting the contraction of R . The contraction-contraction-contraction-contraction-contraction-contraction-contraction-contraction-contraction-contraction-contraction-contraction-contraction-contraction-contractio
    ReportingTime-
 JulianDateTimeRecordType
    R and R are the message R and R are the R -form of R and R are the R -form of R and R and R are the R -form of R -for
    ReportingTrackInfo-
 TrackInfoType
    reporting the contact of th
    ReportingBarrierData-
 BarrierArrayType
end record;
```
Figure Internal message format example -contd

although it may still be too terse for a non-familiar user. For example, a USR message for the external message from section might be

Peterson AFB 001/09:30 - "Hi" - - -

The design of a USR is left up to the designer of the system employing MTV. There appears to be no established formal or informal method for describing USRs This is presumably because the design of a USR is straightforward, given the structure already established in the INR.

The Logical Representation

In order to specify the contents of messages independent of their format, a fourth representation is introduced: the *logical representation*. The logical representation is a mathematical representation that precisely delineates what messages can be described by MTV It is also the representation in which message constraints are specified. The logical representation does not, however, describe the physical representation of a message.

Like the INR, the logical representation is a structured representation—unlike the $\rm EXR$ and USR which are just strings of bits and characters. In fact, a given INR will likely be very similar in structure to its logical representation. The logical representation, however, is not tied down to the details of a particular programming language, and enables the specification of a message format and constraints independent of the programming language being used to implement it

A logical message is composed of primitive fields and composite fields. A primitive field in a logical message is one of

 \mathbf{S} . Thus can also be variable length character string (optionally bounded in length \mathbf{S} by *maxlen*)

integer (*range)* An unsigned integer, optionally bounded by range

name An instance of a previously defined logical representation.

Composite fields define the structure of a message. They combine fields into a new field with a given structure. A composite field is one of:

- { label₁: field₁,..., label_n: field_n } A labelled n-ary product of fields. This is like a record in Ada
- \mathbb{P}^{conv} is form \mathbb{P}^{conv} and \mathbb{P}^{conv} is a labelled in any sum of forms. This is like a variant record in Ada where the discriminant is implicit and is an enumeration of $label_1$ labels are described to the sum may just be a union of symbolic complement corresponding to an enumeration in Ada
- $\sum_{i=1}^{n}$
- $\mu e \mu$ A sequence of $\mu e \nu a$ (a list). Essentially an array of arbitrary length. This has no immediate counterpart in Ada, although a sequence is easily implemented using a linked list

A logical representation for the example message is in figure 4.

```
message Julian Date Time = {
      Julian Day integer -

      \cdots integer in the set of \mathcal{A}Minute integer -

\}message Confidence = [High, Medium, Low, None]
message ContactData = \{Track Number in the second state of the se
      ForceCode: [ Hostile, Unknown, Friendly ],
     Position: { Latitude: integer,
                    Longitude: integer \}\}message Barrier Segment = \{Detection_Confidence: Confidence,
      Probability of Detect integer -

\}message Example = \{Reporting Location: [Andrews AFB,
                                 Peterson AFB
                                 Wright_Patterson_AFB ],
      Reporting Time: Julian Date Time,
      Reporting Message Text string -

      Reporting_Altitude_or_Track_Confidence: [Altitude: integer,
                                                         Track_Confidence: Confidence ],
      Reporting_Contact_Data: ContactData*
      Reporting_Barrier_Data: Barrier_Segment[8]
\}
```
Figure Logical Representation Example

3 Message Translation

For each of the three primary MTV representations -EXR INR and USR a translation must be defined to and from the logical representation. In describing these translations, however, it will become evident that they perform the dual role of serving as formal specifications for the representation being targeted -EXR INR and USR Indeed the specication for an EXR -or USR) to Logical translation can easily be seen as a specification for a parser of a particular external or user message formation of the contract of the cont

To describe the possible range of translations required of an MTV system, a set of primitive translation functions and *combinators* that build on them is presented for each of the EXR \leftrightarrow Logical and \cup SR \leftrightarrow Logical translations. A *combinator* is a function that takes other translation functions as arguments. The translation for a message is built up by specifying the translations of the primitive fields, then combining those translations using combinators.

Unlike the requirements for EXR and USR translations the requirements for INR transla tions are very simple. Since the two representations have similar structure, all that is required is to be able to specify the following mappings: integer \leftrightarrow integer, string \leftrightarrow String, product \leftrightarrow record, sum \leftrightarrow variant record, array \leftrightarrow array and sequence \leftrightarrow linked list.

In the following, a *pattern* refers to something with the expressive power of regular expressions. For convenience, a single quoted string indicates a bitstring, and a double quoted string indicates a character string

-EXR to Logical

The translation from EXR to Logical is a mapping from bitstrings to logical representations Each of the following functions and combinators takes the indicated arguments and a bitstring to be translated. Since fields can be variable length, each function or combinator returns both the appropriate logical representation, and the number of bits consumed in translation.

The following are the translation functions for primitive fields:

- **ascii**(length): bits \rightarrow string \times integer Convert a fixed length string of bits, length bytes long, into a character string using an ASCII interpretation. The returned length is length of the input string in bits
- **vascii**(delim): bits \rightarrow string \times integer Convert a variable length string of bits into a character string using an ASCII interpretation. The string is delimited with the string specified in the *delim* argument. The returned length is the length of the ASCII encoded string plus the length of the delimiter, in bits.
- binary(length): bits \rightarrow integer \times integer Convert a string of bits of the specified length into an integer using the binary conversion appropriate for the architecture of implementation. The returned length is the same as the input length.
- $\texttt{asc2int}(\texttt{length})\colon \texttt{bits} \to \texttt{integer} \times \texttt{integer}$ Convert a string of bits, interpreted as an ASCII string of digits *length* bytes long, into an integer. The returned length is length of the input string in bits.

In addition, both binary and asc2int may be modified by a combination of the the simple arithmetic operators $+$, $-$, \times and \div to specify offsets and scaling.

The following are the translation combinators for composite fields:

$$
\{spec_1 \dots spec_m\} \colon \, \mathbf{bits} \, \rightarrow \, \{ \mathit{label}_1 \, \, ; \, \mathit{field}_1 \, , \, \dots, \, \mathit{label}_n \, \, ; \, \mathit{field}_n\} \, \times \, \mathbf{integer}
$$

Convert a bitstring into a logical product structure. The *specs* are either a translation specification for a field

$label_i:rel}$: translation

or a delimiter spec

delim

or simply a number of bits to skip

 \mathbf{S}

There must be exactly one translation spec for each element of the logical structure although they need not be in order. The length returned is the total length consumed by the entire record

spec - spec n bits label- -eld- labeln -eldn integer

Convert a bitstring into a logical sum structure. Each of the *specs* is of the form:

$$
pattern \Rightarrow label_i: translation
$$

where *pattern* is a discrimination pattern that determines whether the bitstring represents a value from the *label*_i element of the sum. There must be exactly one spec for each element of the logical sum although they hece itself in order Each discrimination pattern is applied to the bitstring in order until one matches, then the associated transference is applied. For the case of a summer there where the person is omitted, the *translation* is similarly omitted, and the number of bits consumed is the length of the matching pattern The length returned is the length consumed by the applied translation

- translation (length): bits \rightarrow field length \times integer Convert a bitstring into a logical array of remplate velopment \mathcal{I} in the same value of the same as the same \mathcal{I} is the same controller The length returned is the total length consumed by all the translations
- \mathbf{seq} (*translation, end-pattern, delim*): bits \rightarrow *field* \times integer Convert a bitstring into a logical sequence. The *end-pattern* is a pattern to match on the bitstring to determine when the end of the sequence has been reached. The *delim* is an inter-element delimiter. The specified translation is repeated and the results made into a logical eld women which end pattern matches to me the specific the same as the same as the model type of *translation*. The length returned is the total length consumed by all the translations -including delimiters

In addition, a record translation spec may be modified by a distribution operator. A distribution operator is used to convert a record whose fields are each an array of the same length into an array of records, each field of which is now a singleton. In short, it converts a record of arrays into an array of records. This is indicated by prefixing the record spec with array district It is an extra it of specify a distribution to a record all of whose items are the value array of the same length. Similarly, an array of records may be factored into a record of arrays

using array factor

A translation from EXR to Logical for the Example message is presented in figure 5.

-Logical to EXR

The operation necessary for Logical to EXR are then a straightforward variation on the EXR to Logical operations. Note that it is not necessary here to return the additional length information that was necessary in the EXR \rightarrow Logical case for keeping track of how much of the bitstring was consumed by a translation.

- $\arctan(\tan\theta)$: string \rightarrow bits Convert a logical string of the specified length into an ASCII character bitstring
- **vascii**(*detim*): string \rightarrow bits Convert a variable length logical string into a delimited ASCII character bitstring
- **binary** (*length*): integer \rightarrow bits Convert a logical integer into a bitstring of the specified length using the conversion appropriate for the architecture of implementation.
- \min zasc $($ length): integer \rightarrow bits \cup onvert a logical integer into an ASCII character bitstring of the specified length.
- $\{spec_1\dots spec_m\}\colon\{label_1,\dots,label_n:spec1}\}$ ture into a bitstring. The specs are either a translation specification for a field

$label_i:rel}$: translation

Loc = \vert '00' \Rightarrow Andrews AFB, $'10' \Rightarrow$ Peterson AFB, ' $11' \Rightarrow$ Wright Patterson AFB |

 $\text{Date} = \{ \text{ Julian_Day: binary}(9), \text{Hour: binary}(5), \text{ Minute: binary}(6) \}$

Confidence $=$ ["HH" \Rightarrow High, " $MM" \Rightarrow Medium,$ " LL " \Rightarrow Low, "NN" \Rightarrow None] ContactData = { $\text{asc2int}(4)$, \vert "H" \Rightarrow Hostile, " U " \Rightarrow Unknown, " F " \Rightarrow Friendly]. \cdots . \cdots . \cdots , \cdots $\operatorname{asc2int}(3)$ } Barrier_Segments = $\{$ Detection_Confidence: Confidence[8], $\sqrt[n]{n}$ $Probability_of_Detect: \: \mathrm{asc2int}(1)[8]$ } Example $= \{$ Reporting Location: Loc. Reporting Time: Date, Reporting Message Text vascii- g Reporting_Altitude_or_Track_Confidence: $\vert\,\vert 0\mathord{\cdot} 9\vert \Rightarrow \,{\rm Altitude:} \,\, {\rm asc2int}(2) \,\, \mathord{\cap} \,\, 1000,$ \therefore \Rightarrow Track_Confidence: Confidence] Reporting Contact Data seq-ContactData
END " END ". END OF THE CONTRACT CONTRACT OF THE CONTRACT O Reporting_Barrier_Data: array_distr Barrier_Segments }

Figure 5: $EXR \rightarrow$ Logical translation example

or a delimiter spec

delim

or simply a number of bits to skip

 \mathbf{S}

There must be exactly one translation spec for each element of the logical structure although they need not be in order

 $[spec_1 \ldots spec_n]$: [label₁ : Jiela₁ , ..., label_n : Jiela_n] \rightarrow bits Convert a logical sum structure into a bitstring. Each of the specs is of the form:

$$
label_i:! (tag, translation)
$$

The tag is a bitstring to identify this element of the sum, if necessary. There must be exactly one *spec* for each element of the logical sum, although they need not be in cract, refers the case of a sum element where the jecta to chillen the sum and the translation is similarly omitted.

- *transiation* [length]: π eld|length| \rightarrow **bits** Convert a logical array of the specified length into a sidering into head type jeevs is the same as the nead type of translation.
- \mathbf{seq} *translation, detim):* μ *eld* $\;\rightarrow$ **Dits** Convert a logical sequence into a bitstring. The delimited was samed delimited delimited to a same of professor and the control we then the property of p of translation

In addition, both binary and intased may be modified by a combination of the the simple arithmetic operators $+$, $-$, \times and \div to specify offsets and scaling. Further, distributions and factors, as described in the previous section, may also be applied to these translations.

-USR to Logical

The translation for USR to Logical is very similar to the one for EXR to Logical. The most significant difference is that everything in the USR is characters, not bits.

- **ascii**(length): chars \rightarrow string \times integer Convert a fixed length string of characters, length long into a string using an ASCII interpretation The returned length is length of the input string
- **vascii**(delim): chars \rightarrow string \times integer Convert a variable length string of characters into a string using an ASCII interpretation The string is delimited with the string specified in the *delim* argument. The returned length is the length of the string plus the length of the delimiter
- asc2int(length): chars \rightarrow integer \times integer Convert an ASCII string of digits length bytes long, into an integer. The returned length is length of the input string.
- $\{spec_1\dots spec_m\}\colon$ chars \to $\{label_1,$ $,$ $field_1,$ $\dots,$ $label_n$: $field_n\}\times$ $integral$ \subset \subset \subset \subset \subset acter string into a logical product structure. The *specs* are the same as for EXR to Logical, except that the skip is in chars, not bits. There must be exactly one translation spec for each element of the logical structure although they need not be in order. The length returned is the total length consumed by the entire record.

[spec₁ ...spec_n]: chars \rightarrow [label₁ : field₁, ..., label_n : field_n] \times integer Convert a character string into a logical sum structure. Each of the specs is of the form:

$pattern \Rightarrow label_i : translation$

where *pattern* is a discrimination pattern that determines whether the bitstring represents a value from the *label*_i element of the sum. There must be exactly one spec for each element of the logical sum although they head not be in order Each. discrimination pattern is applied to the bitstring in order until one matches, then the associated transfation is applied. For the case of a sum element where the jecta is omitted, the *translation* is similarly omitted. The length returned is the length consumed by the applied translation

- *translation* [length]: chars \rightarrow *field* [length] \times integer Convert a character string into a logical array of length length The eld type -eld is the same as the eld type of translation The length returned is the total length consumed by all the translations
- $\mathbf{seq}(\textit{translation},\textit{end}$ -pattern, delim): chars \rightarrow field $\overline{\times}$ integer Convert a character string into a logical sequence. The *end-pattern* is a pattern to match on the character string to determine when the end of the sequence has been reached. The *delim* is an interelement delimiter. The specified translation is repeated and the results made into a logical sequence until endpattern matches The eld type -eld is the same as the field type of *translation*. The length returned is the total length consumed by all the translations -including delimiters

-Logical to USR

The translation for Logical to USR is very similar to the one for Logical to EXR

- $\arctan(\tan \theta)$: string \rightarrow chars Convert a logical string of the specified length into an ASCII character string
- **vascii**(*detim*): string \rightarrow chars Convert a variable length logical string into a delimited ASCII character string
- \min zasc $($ length): integer \rightarrow chars Convert a logical integer into an ASCII character string of the specified length.
- $\{spec_1\ldots spec_m\}\colon\{label_1\ldots,~label_1\ldots,~label_n:spec_1\ldots spec_m\}$ for Logical to EXR except that the skip is in chars, not bits. There must be exactly one translation spec for each element of the logical structure although they need not be in order
- $[spec_1 \ldots spec_n] \colon [\textit{label}_1; \textit{break}_n, \ldots, \textit{label}_n] \rightarrow \textbf{cars}$ Convert a logical sum structure into a character string. Each of the specs is of the form:

$$
label_i:star} (tag, translation)
$$

The tag is a character string to identify this element of the sum, if necessary. There must be exactly one *spec* for each element of the logical sum, although they need not be in order For the case of a sum element where the first model of a sum the translation of is similarly omitted

- *transiation* [*length*]: *pela[length*] \rightarrow chars Convert a logical array of the specified length into a character string that iteration is the same and the same as the same α is the same as the same of the
- \mathbf{seq} (*translation, delim):* η e*ld* $\;\rightarrow$ **chars** Convert a logical sequence into a character string. The delim is an interelement delimiter The eld type -eld is the same as the eld type of translation

$\overline{4}$ Message Validity

Messages have associated with them a notion of validity There are two aspects to this

- \bullet intra-neid validity. The data in a neid is in the range of valid data for that neid.
- \bullet inter-neid validity. The data in interdependent nelds conform to specified constraints.

Intra- and inter-field validity are specified by constraints in the logical representation. Intrafield validity is specified in the declarations of each logical field. The valid range of data for and intracel constraint is either implicit $\mathbf{A} = \mathbf{A} \mathbf{A}$ sum or a sum explicit -eg the bounds on an integer eld Intereld validity is specied as propositions on the logical specification that must be satisfied.

The details of an invalid message are given by its *validity indicator*, a value that gives diagnostic information about what's wrong with a message. Each message format has a unique validity indicator tailored to the specifics of that format.

Derived from analysis of the specification of possible translations, the various intra-field constraint indicators are as follows

- message Truncated This indicates that at some point in the translation it was deter mined that the bitstring was too short for the specified translation.
- message Toolong This indicates that at the end of the translation for a message the total number of bits consumed was not equal to the length of the message
- U is the distribution of \mathbb{Z} . This indicate the value of delimiter \mathbb{Z} is a variable length string was not found
- S is indicated the maximum variable S variable length string exceeded its optional maximum imum length
- \blacksquare indicates the indicates that in an an ascint translation that it can eld contained an ASC \blacksquare character that was not a digit $0-9$.

Cut Chroning This indicates that an integer was out of its specifical ranger

MissingDelimiter This indicates that a delimiter specied between two elds in a record, or between two elements of a sequence, did not match the contents of the message at that point

Discrimination Failure This indicates that indicates the discrimination pattern for a university matched that

Universities of a sequence This indicates that the end of a sequence was not found.

An inter-field constraint is given a name for reference in the validity indicator. It is a predicate, defined as follows:

• simple comparisons on integer or string terms: $=$, \neq , \lt , $>$, \leq and \geq (for strings, lexicographic ordering is used

- \bullet complex constraints using connectives $\mathbf{and},\mathbf{or},\mathbf{not}$ and implies
- quantifiers \forall and \exists for expressing constraints over arrays and sequences

An integer or string term is

- \bullet a logical neld name
- a simple arithmetic expression on (integer) terms using $+$, $-$, \times and \div
- \bullet a length function for sequences and variable length strings $-$

For example the following expresses the constraint mentioned in the note at the bottom of the EXR description

length-Reporting Contact Data implies

 $\forall x$ in Reporting Barrier Data, x. Detection Confidence = None

$\overline{5}$ Test generation

The MTV generator must be able to generate its own tests. In particular, for each message format, a set of test data is to be supplied by the designer of the format, who also supplies the appropriate validity indicator. The generated test will translate and validate all the test data, indicating any inconsistencies between the specied validity indicator and the actual validity indicator

6 The MTV artifact

The artifact produced by the MTV generator will be a module in the target language with a datatype for the INR, a declaration for its validity indicator and the following function interface:

- \bullet a value function from EAR to INR $\hspace{0.1em}$
- \bullet an \it{image} function INR to EAR $\hspace{0.1cm}$
- \bullet a *check* function that indicates the validity of a given EXR, returning a validity indicator that describes the nature of the problem
- \bullet a value function from USR to INR $\hspace{0.1em}$
- \bullet an \it{image} function INR to USR $\hspace{0.1mm}$
- \bullet a *check* function that indicates the validity of a given USK

The value and *image* functions are derived from the specified translations to and from the logical representation. In particular:

> value-la **. Extra** . The model and the set of \mathbf{L} is \mathbf{L} . In \mathbf{L} is \mathbf{L} in \mathbf{L} is \mathbf{L} is \mathbf{L} in \mathbf{L} is a subsequently in \mathbf{L} value-^x USR INR LogicaltoINR-USRtoLogical-x $\frac{1}{2}$ interesting the state of $\frac{1}{2}$ in $\frac{$

This module will be realized by an Ada package whose PDL is similar to the one given in figure 6.

```
with MTV_Common;
package MTV_instance is
  -- INTERNAL REPRESENTATION
     type INR is 
     subtype EXR is Bitstring
     -- Functions for converting between EXRs and INRs
     \mathbf{A} in Indian \mathbf{A} in Indian I
     \mathbf{I} - in Exr return In
     -- Function for checking the validity of an EXR
     type EXR_Validity_Indicator is ...;
     function Check
ImageIn -
 in EXR return EXRValidityIndicator
  -- USER REPRESENTATION
     subtype USR is VString
     -- Functions for converting between USRs and INRs
     function is a set of the contract of
     function Value
ImageIn -
 in USR return INR
     -- Function for checking the validity of a USR
     type USR_Validity_Indicator is \ldots;
     \blacksquareIn \blacksquare Constraint raised by Image and Value functions
     ConstraintError -
 exception
end MTV_Instance
```
Figure 6: Ada PDL for an MTV instance

A requirement of the value functions is that they raise an exception when the input message violates constraints. In Ada, the Constraint Error exception will be the one raised. The user may then consult the validity indicator returned by the check function for details of which constraint was violated

Finally, it is required that each of the value/image pairs form a retraction pair, i.e. that the following two conditions always hold

value-image-x ^x

village called α / / w whenever check α indicates a valid message